

INTERTEMPORAL MODELS: BASICS OF DYNAMIC PROGRAMMING

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Introduction

DYNAMIC PROGRAMMING

- ❑ Can we represent intertemporal problems **recursively**?
 - ❑ Rather than **sequentially**
- ❑ Benefits
 - ❑ Allows application of series of theorems/results that guarantee a **solution exists in the space of functions**
 - ❑ Allows application of series of theorems/results that help **find solution in the space of functions**
 - ❑ Computational algorithms require it – computers can't handle infinite-dimensional objects!
- ❑ Costs
 - ❑ May rule out some solutions to the original (sequential) problem
 - ❑ Requires (a lot?) more structure on the problem
 - ❑ Sometimes (often?) not obvious how to recast sequential problem as recursive problem
- ❑ Ljungqvist and Sargent (2004, p. 16)

"The art in applying recursive methods is to find a convenient definition of a state. It is often not obvious what the state is, or even whether a finite-dimensional state exists."

DYNAMIC PROGRAMMING

- ❑ Can we represent intertemporal problems **recursively**?
 - ❑ Rather than **sequentially**
- ❑ **Benefits**
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 - ❑ Allows application of series of theorems/results that help **find solution in the space of functions**
 - ❑ Computational algorithms require it – computers can't handle infinite-dimensional objects!
- ❑ **Costs**
 - ❑ May rule out some solutions to the original (sequential) problem
 - ❑ Requires (a lot?) more structure on the problem
 - ❑ Sometimes (often?) not obvious how to recast sequential problem as recursive problem
- ❑ **Start with deterministic case**
 - ❑ (Fairly) straightforward
 - ❑ Stochastic case requires more structure

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FROM SEQUENTIAL TO RECURSIVE

- ❑ Lagrangian of consumer problem, with planning horizon T

$$V^0(a_{-1}, r_0; \cdot) \equiv \max_{\{c_t, a_t\}_{t=0}^{\infty}} \sum_{t=0}^T \beta^t \left[u(c_t) + \lambda_t (y_t + (1 + r_{t-1})a_{t-1} - c_t - a_t) \right]$$
- ❑ **State variables** of consumer problem at beginning of any period s
 - ❑ a_{s-1} (accumulation variable) – **the critical one b/c a_{τ} , $\tau \geq s$, are choices**
 - ❑ r_s (price-taker)
 - ❑ A sufficient summary of the **dynamic position of the environment** in which the consumer operates

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FROM SEQUENTIAL TO RECURSIVE

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- **State variables** of consumer problem at beginning of any period s
 - a_{s-1} (accumulation variable) – the critical one b/c $a_t, \tau \geq s$, are choices
 - r_s (price-taker)
 - A sufficient summary of the **dynamic position of the environment** in which the consumer operates
- Define $V^0(a_{-1}, r_0; \cdot)$ as value function starting from period zero
 - The maximized **value** of the constrained optimization problem
 - As function of period-zero parameters of the problem
- Goal: recast problem of finding optimal **sequence** $\{c_t, a_t\}_{t=0,1,2,\dots,T}$ into problem of finding **functions** $\{V(\cdot)\}_{t=0,1,2,\dots,T}$
 - (Actually, find $V(\cdot)$ along with two other functions)

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FROM SEQUENTIAL TO RECURSIVE

- Write out more explicitly

$$V^0(a_{-1}, r_0; \cdot) \equiv \max_{\{c_0, a_0, c_t, a_t\}_{t=1}^{\infty}} \left\{ u(c_0) + \lambda_0 (y_0 + (1 + r_{-1})a_{-1} - c_0 - a_0) + \sum_{t=1}^T \beta^t \left[u(c_t) + \lambda_t (y_t + (1 + r_{t-1})a_{t-1} - c_t - a_t) \right] \right\}$$

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FROM SEQUENTIAL TO RECURSIVE

□ Write out more explicitly

$$V^0(a_{-1}, r_0; \cdot) \equiv \max_{\{c_0, a_0, c_t, a_t\}_{t=1}^{\infty}} \left\{ u(c_0) + \lambda_0 (y_0 + (1 + r_{-1})a_{-1} - c_0 - a_0) + \sum_{t=1}^T \beta^t [u(c_t) + \lambda_t (y_t + (1 + r_{t-1})a_{t-1} - c_t - a_t)] \right\}$$

↓ Separate terms

$$V^0(a_{-1}, r_0; \cdot) \equiv \max_{c_0, a_0} \{ u(c_0) + \lambda_0 (y_0 + (1 + r_{-1})a_{-1} - c_0 - a_0) \} + \max_{c_0, a_0} \left\{ \max_{\{c_t, a_t\}_{t=1}^{\infty}} \sum_{t=1}^T \beta^t [u(c_t) + \lambda_t (y_t + (1 + r_{t-1})a_{t-1} - c_t - a_t)] \right\}$$

Note the max inside the max

↓ Adjust β factors

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FROM SEQUENTIAL TO RECURSIVE

Adjust β factors

$$\longrightarrow V^0(a_{-1}, r_0; \cdot) \equiv \max_{c_0, a_0} \{ u(c_0) + \lambda_0 (y_0 + (1 + r_{-1})a_{-1} - c_0 - a_0) \} + \beta \cdot \max_{c_0, a_0} \left\{ \max_{\{c_t, a_t\}_{t=1}^{\infty}} \sum_{t=1}^T \beta^{t-1} [u(c_t) + \lambda_t (y_t + (1 + r_{t-1})a_{t-1} - c_t - a_t)] \right\}$$

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FROM SEQUENTIAL TO RECURSIVE

Adjust β factors \longrightarrow

$$V^0(a_{-1}, r_0; \cdot) \equiv \max_{c_0, a_0} \left\{ u(c_0) + \lambda_0 (y_0 + (1 + r_{-1})a_{-1} - c_0 - a_0) \right\}$$

$$+ \beta \cdot \max_{c_0, a_0} \left\{ \max_{\{c_t, a_t\}_{t=1}^{\infty}} \sum_{t=1}^T \beta^{t-1} [u(c_t) + \lambda_t (y_t + (1 + r_{t-1})a_{t-1} - c_t - a_t)] \right\}$$

$V^0(a_{-1}, r_0; \cdot)$ is value function starting from period 0.

Bellman Principle of Optimality: optimal decisions in the initial period induce a future state, from which (future) decisions are optimal (Bellman, 1957)

$\equiv V^1(a_0, r_1; \cdot)$, value function starting from period 1.

The value resulting from optimal decisions starting from period 1.

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FROM SEQUENTIAL TO RECURSIVE

Adjust β factors \longrightarrow

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$$+ \beta \cdot \max_{c_0, a_0} \left\{ \max_{\{c_t, a_t\}_{t=1}^{\infty}} \sum_{t=1}^T \beta^{t-1} [u(c_t) + \lambda_t (y_t + (1 + r_{t-1})a_{t-1} - c_t - a_t)] \right\}$$

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The value resulting from optimal decisions starting from period 1.

Recursive representation of consumer problem

$$V^0(a_{-1}, r_0; \cdot) \equiv \max_{c_0, a_0} \left\{ u(c_0) + \lambda_0 (y_0 + (1 + r_{-1})a_{-1} - c_0 - a_0) + \beta \cdot V^1(a_0, r_1; \cdot) \right\}$$

- **Bellman Equation**
- Can analyze optimization problem for period zero...
 - ...given **Bellman Principle of Optimality** holds
 - (But how do $V^0(\cdot)$ and $V^1(\cdot)$ relate to each other?)

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BELLMAN EQUATION

□ Bellman Equation

$$V^0(a_{-1}, r_0; \cdot) \equiv \max_{c_0, a_0} \left\{ u(c_0) + \lambda_0 (y_0 + (1 + r_{-1})a_{-1} - c_0 - a_0) + \beta \cdot V^1(a_0, r_1; \cdot) \right\}$$

- Starting point for recursive analysis
- Applicable to finite T -period or $T \rightarrow \infty$ problems
- Construction requires identifying **state variables** of optimization problem
- **T -period problem**
 - Solution involves sequence of functions $V^0(\cdot), V^1(\cdot), \dots, V^{T-1}(\cdot), V^T(\cdot)$
 - $V^i(\cdot)$ functions in general will differ – reflecting time until end of planning horizon
 - E.g., maximized value starting from age = 60 different from maximized value starting from age = 30 (intuitively)
- **Infinite-horizon problem**
 - **Deterministic case:** $V(\cdot) \equiv V^i(\cdot) = V^j(\cdot) \forall i, j$ **Stochastic case?**
 - Always an infinity of periods left to go **Requires more structure...**

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BELLMAN EQUATION

□ Bellman Equation (for $T \rightarrow \infty$)

$$V(a_{-1}, r_0; \cdot) \equiv \max_{c_0, a_0} \left\{ u(c_0) + \lambda_0 (y_0 + (1 + r_{-1})a_{-1} - c_0 - a_0) + \beta \cdot V(a_0, r_1; \cdot) \right\}$$

- Use to characterize optimal decisions
- Period-0 FOCs

$$c_0: u'(c_0) - \lambda_0 = 0$$

$$a_0: -\lambda_0 + \beta V_1(a_0, r_1; \cdot) = 0 \quad \text{How to compute } V_1(\cdot)?$$

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BELLMAN EQUATION

□ **Bellman Equation (for $T \rightarrow \infty$)**

$$V(a_{-1}, r_0; \cdot) \equiv \max_{c_0, a_0} \left\{ u(c_0) + \lambda_0 (y_0 + (1 + r_{-1})a_{-1} - c_0 - a_0) + \beta \cdot V(a_0, r_1; \cdot) \right\}$$

□ **Use to characterize optimal decisions**

□ **Period-0 FOCs**

$$c_0: u'(c_0) - \lambda_0 = 0$$

$$a_0: -\lambda_0 + \beta V_1(a_0, r_1; \cdot) = 0 \quad \text{How to compute } V_1(\cdot)?$$

□ **Suppose** optimal choice characterized by $c_0 = c(a_{-1}; \cdot)$, $a_0 = a(a_{-1}; \cdot)$
 (c(.) and a(.) **time-invariant functions** in infinite-period problem)

Return to this ...

□ **Insert in value function (can now drop max operator)**

$$V(a_{-1}, r_0; \cdot) \equiv u(c(a_{-1})) + \lambda_0 (y_0 + (1 + r_{-1})a_{-1} - c(a_{-1}) - a(a_{-1})) + \beta \cdot V(a(a_{-1}), r_1; \cdot)$$

□ **Now compute marginal**

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BELLMAN EQUATION

□ **Bellman Equation (for $T \rightarrow \infty$)**

$$V(a_{-1}, r_0; \cdot) \equiv u(c(a_{-1})) + \lambda_0 (y_0 + (1 + r_{-1})a_{-1} - c(a_{-1}) - a(a_{-1})) + \beta \cdot V(a(a_{-1}), r_1; \cdot)$$

□ **Now compute marginal (suppress r argument of c(.) and a(.) functions)**

$$V_1(a_{-1}, r_0; \cdot) =$$

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BELLMAN EQUATION

□ Bellman Equation (for $T \rightarrow \infty$)

$$V(a_{-1}, r_0; \cdot) \equiv u(c(a_{-1})) + \lambda_0 (y_0 + (1 + r_{-1})a_{-1} - c(a_{-1}) - a(a_{-1})) + \beta \cdot V(a(a_{-1}), r_1; \cdot)$$

- Now compute marginal (suppress r argument of $c(\cdot)$ and $a(\cdot)$ functions)

$$V_1(a_{-1}, r_0; \cdot) =$$

$$\Rightarrow V_1(a_{-1}, r_0; \cdot) =$$

Envelope
Condition

□ Envelope Theorem

Note: envelope theorem has nothing to do with dynamic programming

- In computing first-order effects of changes in a problem's **parameters** on the maximized value, can ignore how optimal choices will adjust
- Intuition: because already at a max (marginal costs = marginal benefits)
- Need only consider the direct effect

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BELLMAN EQUATION

□ Bellman Equation (for $T \rightarrow \infty$)

$$V(a_{-1}, r_0; \cdot) \equiv u(c(a_{-1})) + \lambda_0 (y_0 + (1 + r_{-1})a_{-1} - c(a_{-1}) - a(a_{-1})) + \beta \cdot V(a(a_{-1}), r_1; \cdot)$$

- Use to characterize optimal decisions

- Period-0 FOCs, now evaluated using $c(a_{-1})$, $a(a_{-1})$

$$c_0: u'(c(a_{-1})) - \lambda_0 = 0$$

$$a_0: -\lambda_0 + \beta V_1(a_0(a_{-1}), r_1; \cdot) = 0$$

$$\text{Env: } V_1(a(a_{-1}), r_1; \cdot) = \lambda_1(1 + r_0)$$

$$\left. \begin{array}{l} c_0: u'(c(a_{-1})) - \lambda_0 = 0 \\ a_0: -\lambda_0 + \beta V_1(a_0(a_{-1}), r_1; \cdot) = 0 \end{array} \right\} u'(c(a_{-1})) = \beta(1 + r_0)u'(c(a_0))$$

- Seems like usual Euler equation from sequential analysis (deterministic)...

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DETERMINISTIC – RECURSIVE ANALYSIS

- Solution of infinite-horizon consumer problem (starting from date zero)...
- ...is a consumption **decision rule** $c(a_{-1}; \cdot)$, asset **decision rule** $a(a_{-1}; \cdot)$, and **value function** $V(a_{-1}; \cdot)$ that satisfies

- **Bellman equation**

$$V(a_{-1}, r_0; \cdot) \equiv u(c(a_{-1})) + \lambda_0 (y_0 + (1 + r_{-1})a_{-1} - c(a_{-1}) - a(a_{-1})) + \beta \cdot V(a(a_{-1}), r_1; \cdot)$$

- **Euler equation**

$$u'(c(a_{-1})) = \beta V_1(a(a_{-1}), r_1; \cdot) \quad \xleftrightarrow{\text{by envelope theorem}} \quad u'(c(a_{-1})) = \beta(1 + r_0)u'(c(a_{-1}))$$

- which is the TVC in the limit $t \rightarrow \infty$: $\lim_{t \rightarrow \infty} \beta^t u'(c(a_{t-1}^*)) \cdot a(a_{t-1}^*) = 0$

- **Budget constraint**

$$y_0 + (1 + r_{-1})a_{-1} - c(a_{-1}) - a(a_{-1}) = 0$$

taking as given (a_{-1}, r_0, r_{-1})

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DETERMINISTIC – SEQUENTIAL ANALYSIS

- Solution of infinite-horizon consumer problem (starting from date zero)...
- is a consumption and asset **sequence** $\{c_t^*, a_t^*\}_{t=0}^\infty$ that satisfies

- **Sequence of flow budget constraints**

$$c_t^* + a_t^* = y_t + (1 + r_{t-1})a_{t-1}^*, \quad t = 0, 1, 2, \dots$$

- **Sequence of Euler equations**

$$u'(c_t^*) = \beta u'(c_{t+1}^*)(1 + r_t), \quad t = 0, 1, 2, \dots$$

- which is the TVC in the limit $t \rightarrow \infty$: $\lim_{t \rightarrow \infty} \beta^t u'(c_t^*) a_t^* = 0$

taking as given $(\{r_t, y_t\}_{t=0}^\infty, a_{-1}, r_{-1})$

Does solution to recursive problem coincide with solution to sequential problem?

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RECURSIVE VS. SEQUENTIAL ANALYSIS

- ❑ Does solution to recursive problem coincide with solution to sequential problem?
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RECURSIVE VS. SEQUENTIAL ANALYSIS

- ❑ Does solution to recursive problem coincide with solution to sequential problem?
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- ❑ In constructing Bellman representation ($T \rightarrow \infty$ case), the **imposition of time-invariant functions $c(a)$, $a(a)$** potentially limits the class of solutions
 - ❑ In original sequential formulation, this is neither explicitly nor implicitly a requirement of the solution!
- ❑ In general (here without proof...)
 - ❑ Solution to the sequential problem is also a solution to the recursive problem
 - ❑ Solution to the recursive problem is also a solution to the sequential problem **provided some further regularity conditions hold**
- ❑ Stokey, Lucas, Prescott text (1989)

RECURSIVE VS. SEQUENTIAL ANALYSIS

□ So why go recursive?

- Underlying theory:
Contraction Mapping Theorem, Blackwell's Sufficient Conditions for a Contraction, Theorem of the Maximum
- Allows application of series of theorems/results that guarantee a **solution exists in the space of functions**
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 - Computational algorithms require it – computers can't handle infinite-dimensional objects!
 - Econ 701, 702, 630: various computational algorithms

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RECURSIVE VS. SEQUENTIAL ANALYSIS

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 - Computational algorithms require it – computers can't handle infinite-dimensional objects!
 - Econ 701, 702, 630: various computational algorithms
 - Can't really "choose" whether want to analyze problem sequentially or recursively
 - All but the most limited of problems/models require computational solution
 - In which case model analysis **is** recursive
 - What about **stochastic** dynamic programming?
 - Even more structure required....

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STOCHASTIC DYNAMIC PROGRAMMING

- ❑ Even more structure required on the problem to recursively solve dynamic stochastic optimization problems
- ❑ **Main (new) technical problem**
 - ❑ **Branching** of event tree at each of T periods (possibly $T \rightarrow \infty$)
- ❑ **Main technical solution/assumption**
 - ❑ Assume risk follows **Markov** process
 - ❑ Which enables series of theorems/results from deterministic dynamic programming to work in stochastic case...
 - ❑ ...given further technical regularity assumptions (Econ 602)
- ❑ Illustrate the technical problem by building on the stochastic two-period model