# BASICS OF DYNAMIC PROGRAMMING (CONTINUED)

**JANUARY 26, 2012** 

#### Macro Fundamentals **RECURSIVE REPRESENTATION State variables** A sufficient summary, as of the start of period t, of the dynamic position of the environment in which the maximizing agent operates "Environment" of the agent - what needs to be known in order to optimize in period t? ☐ Individual-specific quantities Important: states can be The usual Market prices endogenous or exogenous suspects Government policies (Fixed structural parameters - will omit from state vector for parsimony) Ljungqvist and Sargent (2004, p. 16) "The art in applying recursive methods is to find a convenient definition of a state. It is often not obvious what the state is, or even whether a finite-dimensional state exists." January 26, 2012

Macro Fundamentals

It, of the dynamic imizing agent operates be known in order to int: states can be lous or exogenous state vector for parsimony) intities, prices, govt optimize in period t

#### RECURSIVE REPRESENTATION

#### State variables

- $\ \square$  A sufficient summary, as of the start of period t, of the dynamic position of the environment in which the maximizing agent operates
- "Environment" of the agent what needs to be known in order to optimize in period t?

The usual suspects

- ☐ Individual-specific quantities ☐ Important: states can be endogenous or exogenous
- □ Government policies
  □ (Fixed structural parameters will omit from state vector for parsimony)
- "Sufficient" there are no other objects (quantities, prices, govt policies, etc.) that must be known in order to optimize in period t
- □ Concept well-defined for both finite-T and  $T \rightarrow \infty$  problems
- ☐ KEY: Period-*t* decisions are function of the period-*t* state variables
- ☐ Ljungqvist and Sargent (2004, p. 16)

"The art in applying recursive methods is to find a convenient definition of a state. It is often not obvious what the state is, or even whether a finite-dimensional state exists."

January 26, 2012

3

Deterministic Dynamic Programming

## **BELLMAN EQUATION**

#### □ Bellman Equation

$$V^{0}(a_{-1}, r_{0}; .) = \max_{c_{0}, a_{0}} \left\{ u(c_{0}) + \lambda_{0} \left( y_{0} + (1 + r_{-1}) a_{-1} - c_{0} - a_{0} \right) + \beta \cdot V^{1}(a_{0}, r_{1}; .) \right\}$$

- □ Starting point for recursive analysis
- □ Applicable to finite *T*-period or  $T \rightarrow \infty$  problems
  - Construction requires identifying state variables

#### □ *T*-period problem

- □ Solution involves sequence of functions  $V^0(.)$ ,  $V^1(.)$ , ...,  $V^{T-1}(.)$ ,  $V^T(.)$
- □ V(.) functions in general will differ reflecting time until end of planning horizon
- E.g., maximized value starting from age = 60 different from maximized value starting from age = 30 (intuitively)

#### ☐ Infinite-horizon problem ("stationary" environment)

- □ Deterministic case:  $V(.) \equiv V'(.) = V(.) \forall i,j$
- ☐ Always an infinity of periods left to go

January 26, 2012

## **BELLMAN EQUATION**

□ Bellman Equation (for  $T \rightarrow \infty$ )

$$V(a_{-1}, r_0; .) = \max_{c_0, a_0} \left\{ u(c_0) + \lambda_0 \left( y_0 + (1 + r_{-1}) a_{-1} - c_0 - a_0 \right) + \beta \cdot V(a_0, r_1; .) \right\}$$

- ☐ Use to characterize optimal decisions
- Period-0 FOCs, evaluated using time-invariant  $c(a_{-1})$ ,  $a(a_{-1})$

$$c_{o}: \quad u'(c(a_{-1})) - \lambda_{0} = 0$$
 
$$a_{o}: \quad -\lambda_{0} + \beta V_{1}(a_{0}(a_{-1}), r_{1}; .) = 0$$
 
$$e \quad u'(c(a_{-1})) = \beta(1 + r_{0})u'(c(a_{0}))$$
 
$$e \quad v_{1}(a(a_{-1}), r_{1}; .) = \lambda_{1}(1 + r_{0})$$

□ Seems like usual Euler equation from sequential analysis (deterministic)...

January 26, 2012

5

Notation

## **BELLMAN EQUATION**

□ Bellman Equation (for  $T \rightarrow \infty$ )

$$V(a_{-1}, r_0; .) \equiv u(c(a_{-1})) + \lambda_0 \left( y_0 + (1 + r_{-1})a_{-1} - c(a_{-1}) - a(a_{-1}) \right) + \beta \cdot V(a(a_{-1}), r_1; .)$$

- □ Seems like a two-period problem
  - ☐ In terms of (value) functions, not in terms of choice variables
  - □ Optimize in current period
  - □ Optimize next period (Bellman's Principle of Optimality)

January 26, 2012

Notation

## **BELLMAN EQUATION**

□ Bellman Equation (for  $T \rightarrow \infty$ )

$$V(a_{-1}, r_0; .) \equiv u(c(a_{-1})) + \lambda_0 \left( y_0 + (1 + r_{-1})a_{-1} - c(a_{-1}) - a(a_{-1}) \right) + \beta \cdot V(a(a_{-1}), r_1; .)$$

- Seems like a two-period problem
  - ☐ In terms of (value) functions, not in terms of choice variables
  - Optimize in current period
  - Optimize next period (Bellman's Principle of Optimality)
- □ Common notation
  - Use x for current-period variables
  - Use x' for next-period variables
- □ Bellman Equation

$$V(a,r;.) \equiv u(\underbrace{c(a)}) + \lambda \left( y + (1+r_{-1})a - \underbrace{c(a)} - \underbrace{a(a)} \right) + \beta \cdot V(\underbrace{a(a)}, r';.)$$

$$= c \qquad = a' \qquad = a'$$

☐ Euler equation

January 26, 2012

7

Dynamic Programming

## **RECURSIVE VS. SEQUENTIAL ANALYSIS**

- □ So why go recursive?
  - Allows application of series of theorems/results that guarantee a solution exists in the space of functions

Allows application of series of theorems/results that help find solution in the space of functions

Theory:
Contraction Mapping Theorem, Blackwell's Sufficient Conditions for a Contraction, Theorem of the Maximum

- □ Suppose *V*(.) exists
- $\square$  Procedure for finding V(.) and associated decision rules: iterate on Bellman Equation starting from any arbitrary initial guess

January 26, 2012

Dynamic Programming

## **RECURSIVE VS. SEQUENTIAL ANALYSIS**

- So why go recursive?
  - Allows application of series of theorems/results that guarantee a solution exists in the space of functions
    - Allows application of series of theorems/results that help find solution in the space of functions

Theory: Contraction Mapping Theorem, Blackwell's Sufficient Conditions for a Contraction, Theorem of the Maximum

Suppose V(.) exists

Underlying

Procedure for finding V(.) and associated decision rules. Equation starting from any arbitrary initial guess – call it V(.) initial guess (some parametric form) Procedure for finding V(.) and associated decision rules: iterate on Bellman

$$V(a,r;.) \equiv \max_{c,a'} \left\{ u(c) + \lambda \left( y + (1+r)a - c - a' \right) + \beta \cdot V^{1}(a',r';.) \right\}$$

January 26, 2012

Dynamic Programming

## **RECURSIVE VS. SEQUENTIAL ANALYSIS**

- So why go recursive?
  - Allows application of series of theorems/results that guarantee a solution exists in the space of functions
- Allows application of series of theorems/results that help find solution in the space of functions Underlying

Theory:

Contraction Mapping Theorem, Blackwell's Sufficient Conditions for a Contraction, Theorem of the Maximum

- Suppose V(.) exists
- Procedure for finding V(.) and associated decision rules: iterate on Bellman Equation starting from any arbitrary initial guess – call it  $V^1(.)$

initial guess (some parametric form)

$$V(a,r;.) \equiv \max_{c,a'} \left\{ u(c) + \lambda \left( y + (1+r)a - c - a' \right) + \beta \cdot V^{1}(a',r';.) \right\}$$

$$\bigcirc \quad \text{Conduct maximization}$$

- - $\Box$  Gives functions c(a) and a(a)
  - These are candidate (optimal) decision rules
- Insert candidate c(a) and a(a) into RHS of Bellman Equation generates  $V^2(.)$

Dynamic Programming

#### **RECURSIVE VS. SEQUENTIAL ANALYSIS**

- ☐ So why go recursive?
  - Allows application of series of theorems/results that guarantee a solution exists in the space of functions
- Allows application of series of theorems/results that help find solution in the space of functions

Contraction Mapping Theorem, Blackwell's Sufficient Conditions for a Contraction, Theorem of the Maximum

- □ Suppose V(.) exists
- $\square$  Procedure for finding V(.) and associated decision rules: iterate on Bellman Equation starting from any arbitrary initial guess call it  $V^1(.)$

initial guess (som parametric form)

$$V(a,r;.) = \max_{c,a'} \left\{ u(c) + \lambda \left( y + (1+r)a - c - a' \right) + \beta \cdot V^{1}(a',r';.) \right\}$$

- 2 Conduct maximization
  - $\Box$  Gives functions c(a) and a(a)
  - ☐ These are candidate (optimal) decision rules
  - Insert candidate c(a) and a(a) into RHS of Bellman Equation generates  $V^2(.)$

Does  $V^2(.) = V^1(.)$ ?

January 26, 2012

Dynamic Programming

11

## RECURSIVE VS. SEQUENTIAL ANALYSIS

- □ So why go recursive?
  - Allows application of series of theorems/results that guarantee a solution exists in the space of functions

 $\begin{tabular}{ll} $\square$ & Allows application of series of theorems/results that help find solution in the space of functions \\ \end{tabular}$ 

Theory:

Contraction Mapping Theorem, Blackwell's Sufficient Conditions for a Contraction, Theorem of the Maximum

- □ Suppose *V*(.) exists
- $\square$  Procedure for finding V(.) and associated decision rules: iterate on Bellman Equation starting from any arbitrary initial guess call it  $V^1(.)$

initial guess (some parametric form)

 $V(a,r;.) = \max_{c,a'} \left\{ u(c) + \lambda \left( y + (1+r)a - c - a' \right) + \beta \cdot V^{1}(a',r';.) \right\}$ Conduct maximization

If no, insert  $V^2(.)$  on RHS and repeat

- ☐ These are candidate (optimal) decision rules

Insert candidate c(a) and a(a) into RHS of Bellman Equation – generates  $V^2(.)$ 

Does  $V^2(.) = V^1(.)$ ? If yes, stop. Have found  $V(.) (= V^2(.) = V^1(.))$ 

January 26, 2012 12

Dynamic Programming

## **RECURSIVE VS. SEQUENTIAL ANALYSIS**

- So why go recursive?
- Allows application of series of theorems/results that guarantee a solution exists in the space of functions
- Allows application of series of theorems/results that help find solution Underlying in the space of functions

Contraction Mapping Theorem, Blackwell's Sufficient Conditions for a Contraction, Theorem of the Maximum

e.g., value iteration

- Computational algorithms require it computers can't handle infinitedimensional objects!
  - Soon: simple computational algorithms
- Can't "choose" whether to analyze problem sequentially or recursively
  - All but the most limited of problems require computational solution
  - In which case model analysis is recursive

January 26, 2012

Dynamic Programming

## **RECURSIVE VS. SEQUENTIAL ANALYSIS**

- So why go recursive?
  - Allows application of series of theorems/results that guarantee a solution exists in the space of functions

Allows application of series of theorems/results that help find solution in the space of functions Underlying

Theory:
Contraction Mapping Theorem, Blackwell's Sufficient Conditions for a Contraction, Theorem of the Maximum

e.g., value function iteration

- Computational algorithms require it computers can't handle infinitedimensional objects!
  - Soon: simple computational algorithms
- Can't "choose" whether to analyze problem sequentially or recursively
  - All but the most limited of problems require computational solution
    - In which case model analysis is recursive
- "Solving model sequentially"
  - $u(c_t) = \beta(1+r_t)u'(c_{t+1})$ Doesn't seem recursive...
    - Imposing recursivity on solution ...but computational implementation requires time-invariant decision rule

 $u(c(a_{t-1})) = \beta(1+r_t)u'(c(a_t))$ 

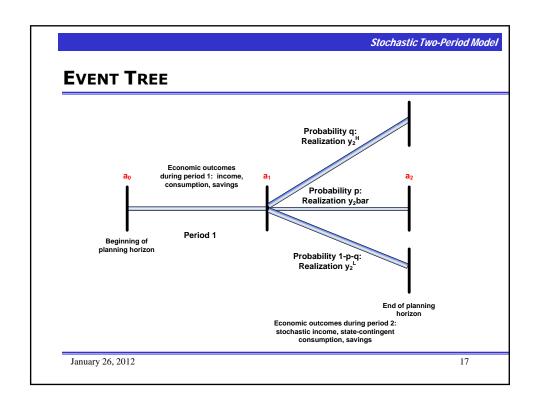
|                                      |   | Dynamic Programming  |  |
|--------------------------------------|---|--|--|
| RE                                   | CUI   | RSIVE VS. SEQUENTIAL ANALYSIS  |  |
| Underlying                           | So why go recursive?  |  |  |
|                                      |   | Allows application of series of theorems/results that guarantee a solution exists in the space of functions Allows application of series of theorems/results that help find solution in the space of functions |  |
| Theory:<br>Contraction               | Mappin  | g Theorem, Blackwell's Sufficient Conditions for a Contraction, Theorem of the Maximum   |  |
| e.g., value<br>function<br>iteration |   | Computational algorithms require it – computers can't handle infinite-dimensional objects!  □ Soon: simple computational algorithms  |  |
|                                      | Can't "choose" whether to analyze problem sequentially or recursively |  |  |
|                                      |   | All but the most limited of problems require computational solution In which case model analysis is recursive  |  |
|                                      | What about stochastic dynamic programming?                            |  |  |
|                                      |   | Even more structure required   |  |
|                                      |   | The key assumption is Markov risk  |  |
| Janu                                 | ary 26,   | 2012 15  |  |
|                                      |   |  |  |

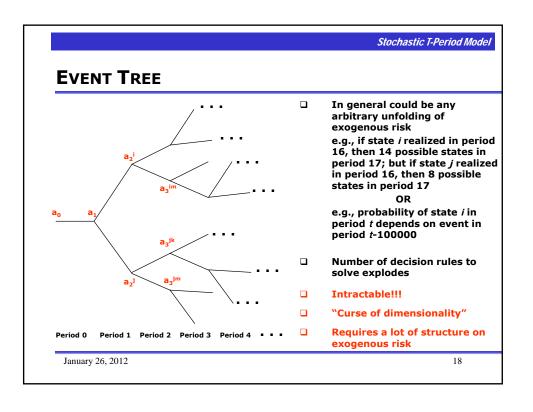
Introduction

### STOCHASTIC DYNAMIC PROGRAMMING

- □ Even more structure required on the problem to recursively solve dynamic stochastic optimization problems
- ☐ Main (new) technical problem
  - □ Branching of event tree at each of T periods (possibly  $T \rightarrow \infty$ )
- □ Main technical solution/assumption
  - ☐ Assume risk follows Markov process
  - Which enables series of theorems/results from deterministic dynamic programming to work in stochastic case...
  - ...given further technical regularity assumptions

January 26, 2012





Macro Fundamentals

#### **RISK STRUCTURE**

#### **Assumptions**

☐ Set of realizations of exogenous state variable is independent of date

$$S2 = S3 = S4 = S5 = .... = ST-1 = ST$$

January 26, 2012

19

Macro Fundamentals

### **RISK STRUCTURE**

#### **Assumptions**

- ☐ Set of realizations of exogenous state variable is independent of date
- Probability of realization of exogenous state variable in period t depends only on outcomes in period t-1
  - Suppose  $X_t$  is a stochastic process and  $X_t$  is a particular realization
  - $X_t$  is a Markov process if

$$\begin{split} \Pr \left( {{X_t} = {x_t}\left| {\left| {{X_{t - 1}} = {x_{t - 1}},{X_{t - 2}} = {x_{t - 2}},{X_{t - 3}} = {x_{t - 3}},....,{X_{t - 10000}} = {x_{t - 10000}},....} \right)} \\ = \Pr \left( {{X_t} = {x_t}\left| {\left| {{X_{t - 1}} = {x_{t - 1}}} \right.} \right.} \right) \\ \text{ CONDITIONAL probability depends on only } \text{ } t\text{-}1} \\ \end{split}$$

January 26, 2012

#### **RISK STRUCTURE**

#### **Assumptions**

- Set of realizations of exogenous state variable is independent of date
- Probability of realization of exogenous state variable in period t depends only on outcomes in period t-1
  - □ Suppose  $X_t$  is a stochastic process and  $X_t$  is a particular realization
  - $\square$   $X_t$  is a Markov process if

$$\begin{split} \Pr \big( X_{t} = x_{t} \mid X_{t-1} = x_{t-1}, X_{t-2} = x_{t-2}, X_{t-3} = x_{t-3}, ...., X_{t-10000} = x_{t-10000}, .... \big) \\ = \Pr \big( X_{t} = x_{t} \mid X_{t-1} = x_{t-1} \big) & \text{CONDITIONAL probability depends on only $t$-1} \end{split}$$

- Not as restrictive as it may seem could have finite lags in process
- □ E.g.
- Just can't have infinite lags (in principle) or "too many" (finite) lags (in computational practice)

January 26, 2012 21

Macro Fundamentals

#### **RISK STRUCTURE**

#### **Assumptions**

- ☐ Set of realizations of exogenous state variable is independent of date
- Probability of realization of exogenous state variable in period t depends only on outcomes in period t-1
  - Suppose  $X_i$ , is a stochastic process and  $X_i$  is a particular realization
  - $X_t$  is a Markov process if

$$\begin{split} \Pr \left( {{X_t} = {x_t}\left| {\left| {{X_{t - 1}} = {x_{t - 1}},{X_{t - 2}} = {x_{t - 2}},{X_{t - 3}} = {x_{t - 3}},....,{X_{t - 10000}} = {x_{t - 10000}},....} \right)} \\ = \Pr \left( {{X_t} = {x_t}\left| {\left| {{X_{t - 1}} = {x_{t - 1}}} \right|} \right.} \right) \\ \text{CONDITIONAL probability depends on only $t$-1} \end{split}$$

- ☐ Not as restrictive as it may seem could have finite lags in process
- □ Exogenous state variable is Markov process + assumption/result that decision rules are time-invariant (for  $T \rightarrow \infty$ ) functions of state variables
  - ⇒ Endogenous processes are Markov given several regularity assumptions

Underlying theory: Stokey, Lucas, Prescott (1989, Chapters 8-12)

January 26, 2012

# STOCHASTIC - SEQUENTIAL ANALYSIS

- Planning horizon  $T \rightarrow \infty$
- Exogenous state drawn from set S (could be continuous or discrete)
- Suppose single asset with state-contingent r (will illustrate main ideas)

$$\max_{\{c_t,a_t\}_{t=0}^T} E_0 \sum_{t=0}^T \beta^t u(c_t) \text{ subject to } \begin{cases} c_t + a_t = y_t + (1+r_t)a_{t-1}, & t = 0,1,2,...T \\ \text{ state-contingent budget constraints in } t > 0 \end{cases}$$

**FOCs** 

$$c_0$$
:  $u'(c_0) - \lambda_0 = 0$ 

 $a_0$ :

$$c_1$$
:  $\beta E_0 u'(c_1) - \beta E_0 \lambda_1 = 0$ 

Holds for each

$$E_0 u'(c_1^j) = E_0 \lambda_1^j, \ \forall j \in S$$

January 26, 2012

Model Analysis

## STOCHASTIC - SEQUENTIAL ANALYSIS

- Planning horizon  $T \rightarrow \infty$
- Exogenous state drawn from set S (could be continuous or discrete)
- Suppose single asset with state-contingent r (will illustrate main ideas)

$$\max_{\{c_i,a_i\}_{t=0}^T} E_0 \sum_{t=0}^T \beta^t u(c_t) \text{ subject to } \begin{cases} c_t + a_t = y_t + (1+r_t)a_{t-1}, & t = 0,1,2,...T \\ \text{state-contingent budget constraints in } t > 0 \end{cases}$$

FOCs

$$c_0$$
:  $u'(c_0) - \lambda_0 = 0$ 

$$\mathbf{a_0:} \quad -\lambda_0 + \beta E_0 \left[ \lambda_1 (1+r_1) \right] = 0$$

$$1 = E_0 \left[ \frac{\beta \lambda_1}{\lambda_0} (1 + r_1) \right]$$

$$c_1$$
:  $\beta E_0 u'(c_1) - \beta E_0 \lambda_1 = 0$ 

Holds for each

$$E_0 u'(c_1^j) = E_0 \lambda_1^j, \ \forall j \in S$$

# STOCHASTIC - SEQUENTIAL ANALYSIS

- Planning horizon  $T \rightarrow \infty$
- Exogenous state drawn from set S (could be continuous or discrete)
- Suppose single asset with state-contingent r (will illustrate main ideas)  $\max_{\{c_t,a_t\}_{t=0}^T} E_0 \sum_{t=0}^T \beta^t u(c_t) \ \ \text{subject to} \left\{ \begin{array}{l} c_t + a_t = y_t + (1+r_t)a_{t-1}, \quad t = 0,1,2,...T \\ \text{state-contingent budget constraints in } t > 0 \end{array} \right.$

**FOCs** 

$$c_1$$
:  $\beta E_0 u'(c_1) - \beta E_0 \lambda_1 = 0$ 

Holds for each state 
$$E_0 u \, {}^{\text{\tiny $(C_1^j)$}} = E_0 \lambda_1^j, \ \, \forall j \in S$$

**a<sub>1</sub>:** 
$$-\beta E_0 \lambda_1 + \beta^2 E_0 [\lambda_2 (1 + r_2)] = 0$$

$$c_2$$
:  $\beta^2 E_0 u'(c_2) - \beta^2 E_0 \lambda_2 = 0$ 

Holds for each state 
$$E_0 u \, {}^{\text{\tiny $}}(c_2^{\, j}) = E_0 \lambda_2^{\, j}, \ \, \forall j \in S$$

January 26, 2012

Model Analysis

#### STOCHASTIC - MARKOV SOLUTION

- $\{X_t\}_{t=0,1,2,\dots}$  is Markov process (exogenous and endogenous states)
  - Nothing about the probability distribution of  $X_{t+2}$  is known in period tthat is not known in period t+1
    - □ Information set of period t+1 is superset of information set of period t
- Allows applying a law of iterated expectations

$$\Box$$
  $E_{t} X_{t+2} = E_{t} [E_{t+1} X_{t+2}]$ 

$$E_0\lambda_1 = \beta {\color{red} E_0} \Big[ \lambda_2 (1+r_2) \Big] \hspace{1cm} \longleftarrow \hspace{1cm} E_0\lambda_1 = \beta {\color{red} E_0} \Big[ {\color{red} E_1} \Big( \lambda_2 (1+r_2) \Big) \Big]$$

#### STOCHASTIC - MARKOV SOLUTION

- $\{X_t\}_{t=0,1,2,\dots}$  is Markov process (exogenous and endogenous states)
  - Nothing about the probability distribution of  $X_{t+2}$  is known in period tthat is not known in period t+1
    - Information set of period t+1 is superset of information set of period t
- Allows applying a law of iterated expectations
  - $E_{t} X_{t+2} = E_{t} [E_{t+1} X_{t+2}]$

$$E_0\lambda_1 = \beta E_0 \left[\lambda_2(1+r_2)\right] \quad \longrightarrow \quad E_0\lambda_1 = \beta E_0 \left[E_1 \left(\lambda_2(1+r_2)\right)\right] \quad \longleftarrow$$

Date- and state-contingent decisions: decisions governed by this Euler condition are conditional on information set of period 1 (i.e., recursivity) 

$$E_1 \lambda_1 = \beta E_1 \left[ \lambda_2 (1 + r_2) \right] \qquad \longrightarrow \qquad \lambda_1 = \beta E_1 \left[ \lambda_2 (1 + r_2) \right]$$

January 26, 2012

Model Analysis

# STOCHASTIC - SEQUENTIAL ANALYSIS

- Planning horizon  $T \rightarrow \infty$
- Exogenous state drawn from set S (could be continuous or discrete)
- Suppose single asset with state-contingent r (will illustrate main ideas)

$$\max_{\{c_t,a_t\}_{t=0}^T} E_0 \sum_{t=0}^T \beta^t u(c_t) \text{ subject to } \begin{cases} c_t + a_t = y_t + (1+r_t)a_{t-1}, & t = 0,1,2,...T \\ \text{state-contingent budget constraints in } t > \mathbf{0} \end{cases}$$

**FOCs** 

**c<sub>1</sub>:** 
$$\beta E_0 u'(c_1) - \beta E_0 \lambda_1 = 0$$

 $E_1 u'(c_1^j) = E_1 \lambda_1^j, \ \forall j \in S$ 

$$a_1: -\beta E_0 \lambda_1 + \beta^2 E_0 [\lambda_2 (1+r_2)] = 0$$

Because Markov and

 $\lambda_1 = \beta E_1 \left[ \lambda_2 (1 + r_2) \right]$ 

$$c_2$$
:  $\beta^2 E_0 u'(c_2) - \beta^2 E_0 \lambda_2 = 0$ 

Holds for each

 $E_2 u'(c_2^j) = E_2 \lambda_2^j, \ \forall j \in S$ 

#### STOCHASTIC - MARKOV SOLUTION

- Planning horizon  $T \rightarrow \infty$
- Exogenous state drawn from set S (could be continuous or discrete)
- Suppose single asset with state-contingent r (will illustrate main ideas)

$$\max_{\{c_t,a_t\}_{t=0}^T} E_0 \sum_{t=0}^T \beta^t u(c_t) \text{ subject to } \begin{cases} c_t + a_t = y_t + (1+r_t)a_{t-1}, & t = 0,1,2,...T \\ \text{with uncertain realizations in } t > \mathbf{0} \end{cases}$$

**FOCs** 

$$c_t: \quad \beta^t E_0 u'(c_t) - \beta^t E_0 \lambda_t = 0$$

Holds for each date and state

$$u'(c_1^j) = \lambda_1^j, \ \forall j \in S$$

$$\textbf{a_t:} \quad -\beta^t E_0 \lambda_t + \beta^{t+1} E_0 \left[ \lambda_{t+1} (1+r_{t+1}) \right] = 0$$
Because Markov and state- and date-

$$\lambda_{t} = \beta E_{t} \left[ \lambda_{t+1} (1 + r_{t+1}) \right]$$

$$c_{t+1}$$
:  $\beta^{t+1}E_0u'(c_{t+1}) - \beta^{t+1}E_0\lambda_{t+1} = 0$ 

Holds for each date and state

$$u'(c_2^j) = \lambda_2^j, \ \forall j \in S$$

One-period-ahead conditional expectation governs stochastic Euler condition

January 26, 2012

29

Model Solution

#### STOCHASTIC - MARKOV SOLUTION

- Denote exogenous state variables as z (e.g.,  $z_t = [y_t, r_t]$ )
- Solution of infinite-horizon consumer problem is a consumption decision rule c(a, z;.), asset decision rule a(a, z;.), and value function V(a, z;.) that satisfies

Model Solution

#### STOCHASTIC - MARKOV SOLUTION

- Denote exogenous state variables as z (e.g.,  $z_t = [y_t, r_t]$ )
- Solution of infinite-horizon consumer problem is a consumption decision rule c(a, z;.), asset decision rule a(a, z;.), and value function V(a, z;.) that satisfies
  - ☐ (Stochastic) Euler equation

$$u'(c(a,z)) = \beta E \left[ u'(c(a',z'))(1+r') \right]$$

which is the (expectational) TVC in the limit  $t \rightarrow \infty$ :

$$\lim E_0 \beta' u'(c(a,z)) \cdot a(a,z) = 0$$

Budget constraint

$$y + (1+r)a - c(a,z) - a(a,z) = 0$$

January 26, 2012

**Model Solution** 

31

#### STOCHASTIC - MARKOV SOLUTION

- Denote exogenous state variables as z (e.g.,  $z_t = [y_t, r_t]$ )
- Solution of infinite-horizon consumer problem is a consumption decision rule c(a, z;.), asset decision rule a(a, z;.), and value function V(a, z;.) that satisfies
  - ☐ (Stochastic) Euler equation

$$u'(c(a,z)) = \beta E[u'(c(a',z'))(1+r')]$$

□ which is the (expectational) TVC in the limit  $t \rightarrow \infty$ :

$$\lim E_0 \beta^t u'(c(a,z)) \cdot a(a,z) = 0$$

□ Budget constraint

$$y + (1+r)a - c(a,z) - a(a,z) = 0$$
 Expectation in Bellman Equation  $z \rightarrow z'$ 

□ Bellman Equation



taking as given (y,a,r) and (Markov) transition function for  $z \rightarrow z'$ 

January 26, 2012 32

Stochastic Dynamic Programming

## **BELLMAN EQUATION**

□ Bellman Equation (for  $T \rightarrow \infty$ )

$$V(a,z;.) \equiv \max_{c,a'} \left\{ u(c) + \lambda \left( y + (1+r)a - c - a' \right) + \beta \cdot \frac{EV(a',z';.)}{\uparrow} \right\}$$

- ☐ Use to characterize optimal decisions Expectation in Bellman Equation 7
- □ Current-period FOCs, evaluated using c(a,z;.), a(a,z;.)

c:

*a*':

Env:

January 26, 2012

Stochastic Dynamic Programming

## **BELLMAN EQUATION**

□ Bellman Equation (for  $T \rightarrow \infty$ )

$$V(a,z;.) = \max_{c,a'} \left\{ u(c) + \lambda \left( y + (1+r)a - c - a' \right) + \beta \cdot \underbrace{EV(a',z';.)}_{A} \right\}$$

- Use to characterize optimal decisions  $\begin{bmatrix} Expectation in Bellman Equation \\ z \rightarrow z' \end{bmatrix}$
- □ Current-period FOCs, evaluated using c(a,z;.), a(a,z;.)

c: 
$$u'(c(a,z)) - \lambda = 0$$
  
a':  $-\lambda + \beta EV_1(a(a,z), z(a,z);.) = 0$   $u'(c(a,z)) = \beta E[u'(c(a,z))(1+r)]$   
Env:  $EV_1(a,z;.) = \lambda(1+r)$ 

- ☐ Bellman analysis goes through as in deterministic case
  - ☐ (Given further technical conditions we won't study see SLP)

January 26, 2012 34

|     | Macro Fundament.   |
|-----|--|
| M   | ARKOV RISK   |
|     | Why does Markov assumption make everything work?   |
|     | Main issue in moving from deterministic dynamic programming to stochastic dynamic programming: preserving recursivity                                  |
|     | ☐ So exogenous states must also have recursive structure   |
|     | Shocks that have this recursive structure are Markov processes   |
|     | Markov has property that given the current realization, future realizations are independent of the past  "Limited history dependence"  "Finite memory" |
|     | In environments in which the "regularity conditions" that ensure<br>standard Bellman analysis applies to stochastic problems are not<br>satisfied      |
|     | often simply need to ASSUME decision rules are Markov to mak progress  |
| Ian | nary 26, 2012 35   |

